Synchome

Evaluation of Synchronization Application for Mobile Content Retrieval

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Roadmap

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- > What is CCNx
- > CCNx synchronization protocol
- > Problem statement
- > Synchome design
- > Evaluation

What is CCNx

- > CCN = Content Centered Networking
- Communication based on names
- Everything is a message Interest and Data messages
- > CCNx = open source implementation of CCN

CCNx synchronization protocol

- Keeps data automatically synchronized
- > Data is stored in persistent storage repositories
- > User defines what to synchronize
- > New content is synchronized between participating hosts

CCNx synchronization protocol

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- Synchronizing hosts must be in same local subnet or connected by tunnel → additional overhead for mobile clients
- Synchronization runs infinitely unless aborted by user → continued energy expenditure
- Synchronization protocol continues exchanging messages when not actively synchronizing content → inefficient for mobile clients





- > Synchronization protocol for mobile clients
- Registers hosts IP addresses as locators in CCNx routing tables
- Detects synchronization disrupts (e.g. no AP in vicinity) and re-registers hosts if necessary
- > Exits once synchronization of content is finished

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Synchome

- Consists of:
- Client application (registers server, sends initialization Interest to server, runs monitoring application, exits once finished succesfully)
- Server application (home repository, constantly running and listening for Interests)
- > Synchronization monitoring application



Evaluation

- > synchome vs. ccnsyncslice (implementation of synchronization protocol)
- Synchronization of content with/without periods where no new content is available
- Measure synchronization time
- > Use CCNx log to analyze Interest and Data messages passed between client and host

Evaluation



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> Evaluated on UBELIX using DCE-ns3





- Keeps information stored in synchronization tree. Parent folder hash is tree roothash, subfolder hashes are child nodes, etc
- Rootadvise carrying roothash sent every 20 seconds. If roothashes differ, new content available
- When synchronizing, rootadvise Interests are exchanged every second
- Node fetch Interests carrying child node hashes exchanged simultaneously with content until synchronization finished

CCNx synchronization protocol



Scenario 1

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- Initialize CCNx and repositories
- > Add files sized 1, 5 or 10MB to repository
- > 5 files each added to 1, 2 or 3 collections
- Run synchome/ccnsyncslice until finished synchronization and stop simulation

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2 collections, 1-5-10MB





3 collections, 10MB



- > Synchronize for 24 hours
- > Add new content (filesize 1 or 2MB) every 900, 1800 or 3600 seconds
- Consyncesice constantly running
- Synchome is initialized when new content is available and exits when synchronization is finished

Message types



- Interest_to messages: Interests sent out by host. May be Rootadvise or Node fetch Interests, or requests for content.
- Content_to messages: Data messages sent out by host. May carry replies to Rootadvise or Node fetch Interests, or content.





Client Data messages







- Ccnsyncslice sends out around 50% more Data messages than synchome
- Number of transferred content segments stays constantly in same range
- So what is causing the increase?

Client Rootadvise replies





Client Node fetch replies



Conclusions

- Synchome generates far less messages than ccnsyncslice.
- The larger the content generation interval, the better synchome performs compared to consyncalice
- Exchange of Node fetch Interests has influence on synchronization time
- Mean runtime difference is nearly negligible between applications. Ccnsyncslice always displays higher maximal runtime.

Questions?

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