

**PhD defense, June 16, 2005**  
**Addressing the Challenges for TCP over  
Multihop Wireless Networks**

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# Outline

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- > Motivation
- > Multihop networks
- > Transmission Control Protocol (TCP)
- > TCP challenges and proposed solutions
- > ***Packet loss discrimination using Fuzzy Logic***
- > ***Smart acknowledgment strategy***
- > Conclusions and outlook

# Motivation

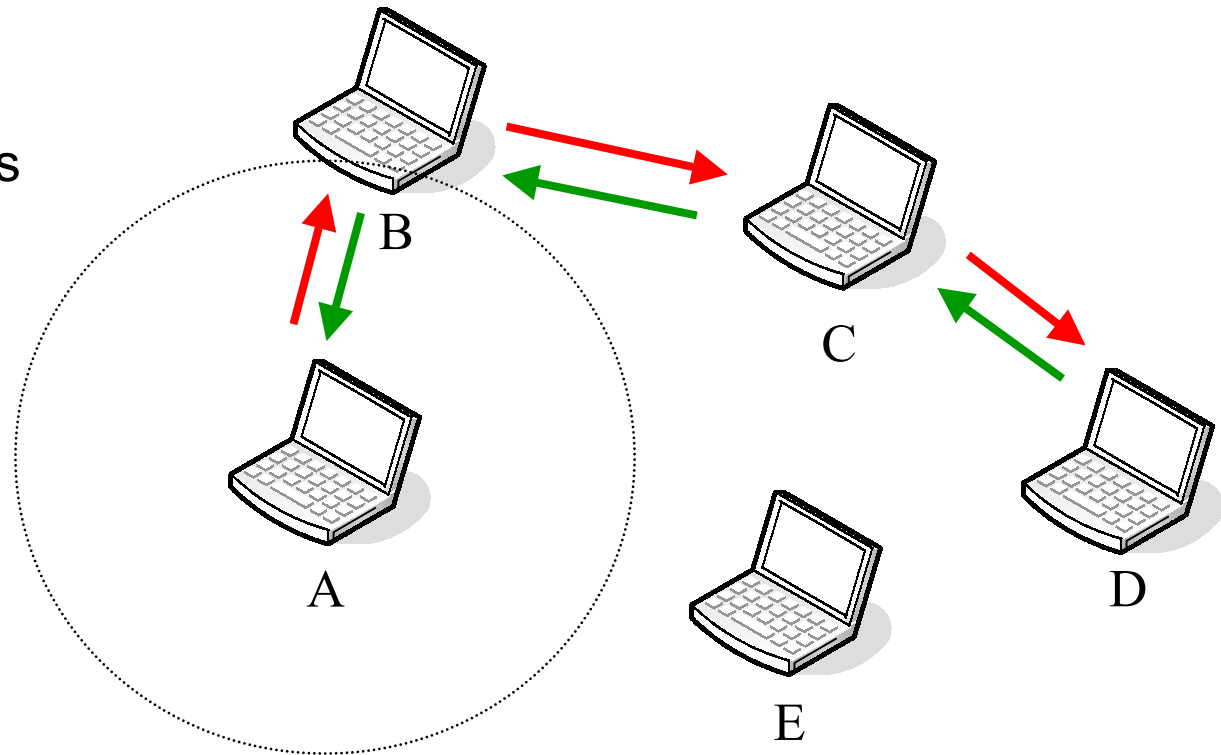
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- > More and more Internet traffic are foreseen in future wireless communications based on IEEE 802.11 standard
  
- > Wireless communications are quite different from the wired world
  - Wireless medium are inherently noisy
  - Bandwidth in wireless channels is a scarce resource
  
- > The widespread use of TCP in the Internet motivates its extension to wireless

# Multihop Wireless Networks

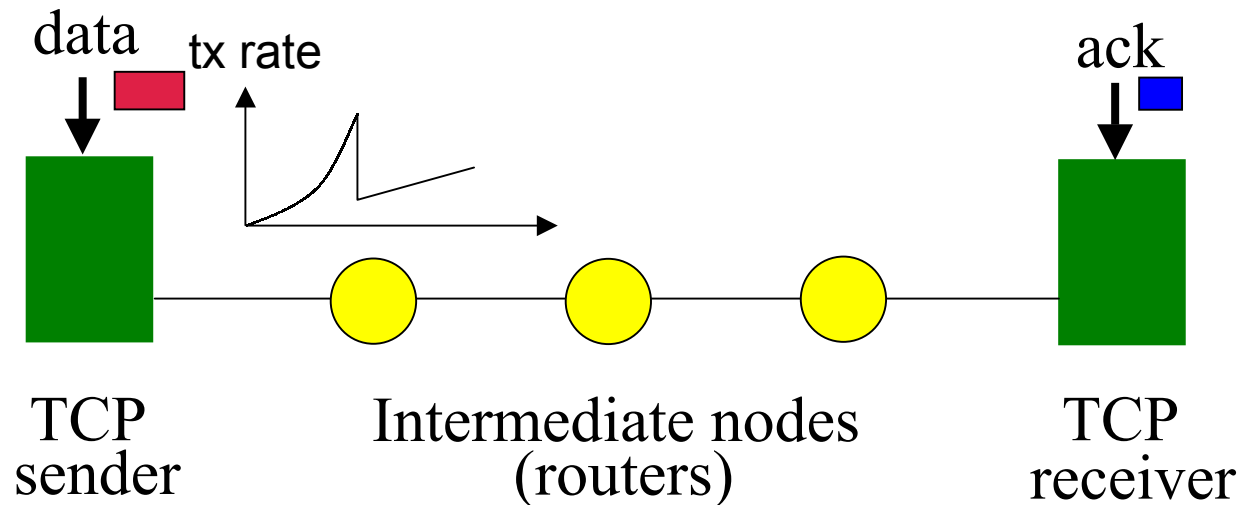
## Multihop ad hoc network

- > No fixed infrastructure
- > Communication range is increased by multiple links
- > Nodes forward data to each other
- > Nodes may move
- > Low energy resources



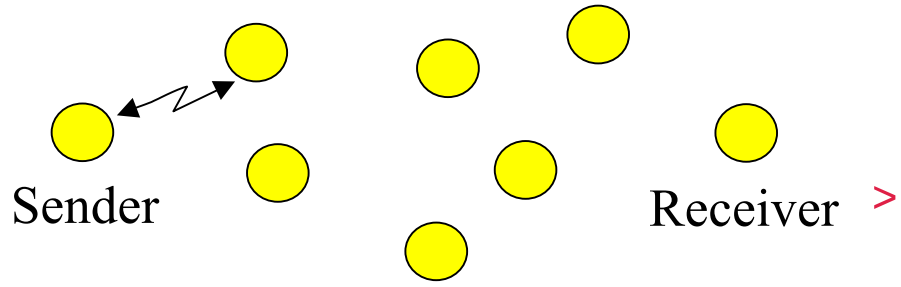
# Transmission Control Protocol (TCP)

- > Provides reliable data delivery (bidirectional flow)
- > Works on an end-to-end basis
- > Continuously probes the network for resources
- > Retransmission: - *retransmit timeout (timer expiration)*  
- *fast retransmit (3 duplicate ACKs)*

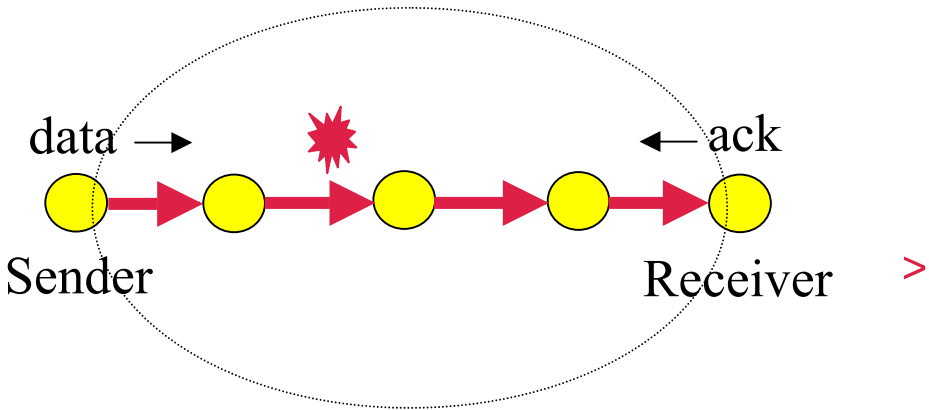


# Challenges and Solutions for TCP

## Wireless network



Lossy channel  
↓  
Packet loss discrimination



Traffic redundancy  
↓  
Smart ACK management

# First Main Contribution: Packet Loss Discrimination

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- > Purpose:
  - To distinguish between congestion and medium induced errors
  
- > How:
  - Evaluating packet delay variations (Round-Trip Time) inside the network
  
- > Advantages:
  - Very lightweight (only a few logic operations)
  - Independent of explicit signaling from the network (end-to-end)
  
- > Assumptions:
  - Fixed packet size
  - Conventional lower layer protocols
  - Known RTT distributions as a function of the # of hops

# Design Rationale

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## Why RTT measurements?

- > TCP relies naturally on round-trip time measurements for taking decisions
- > It is intuitive that such measurements may be useful in inferring the internal network state

## Difficulties

- > RTT measurements may contain overlapping values under certain conditions

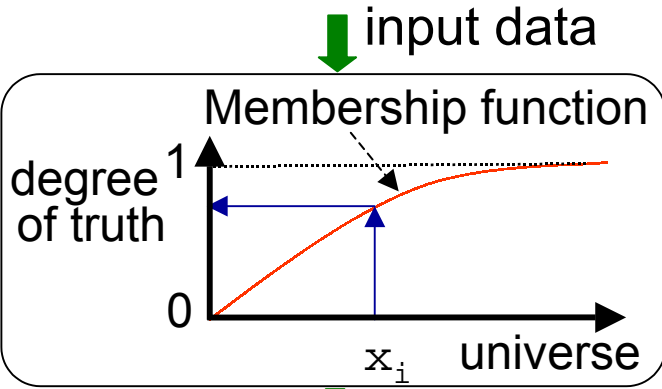
## Then

- > Appropriate tool may perform good reasoning on the measured data → Fuzzy logic

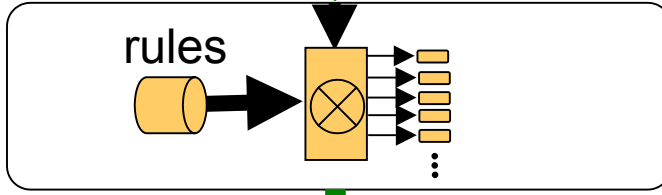


# Fuzzy Logic (FL)

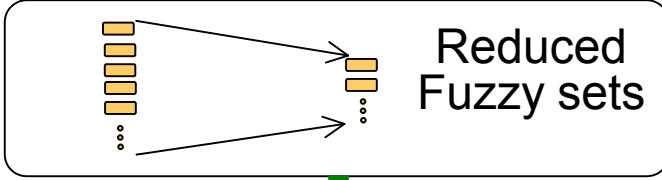
Fuzzification



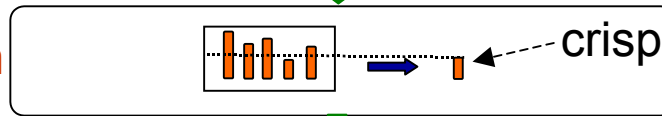
Inference



Aggregation



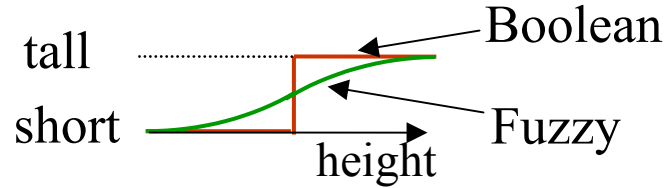
Defuzzification



output data

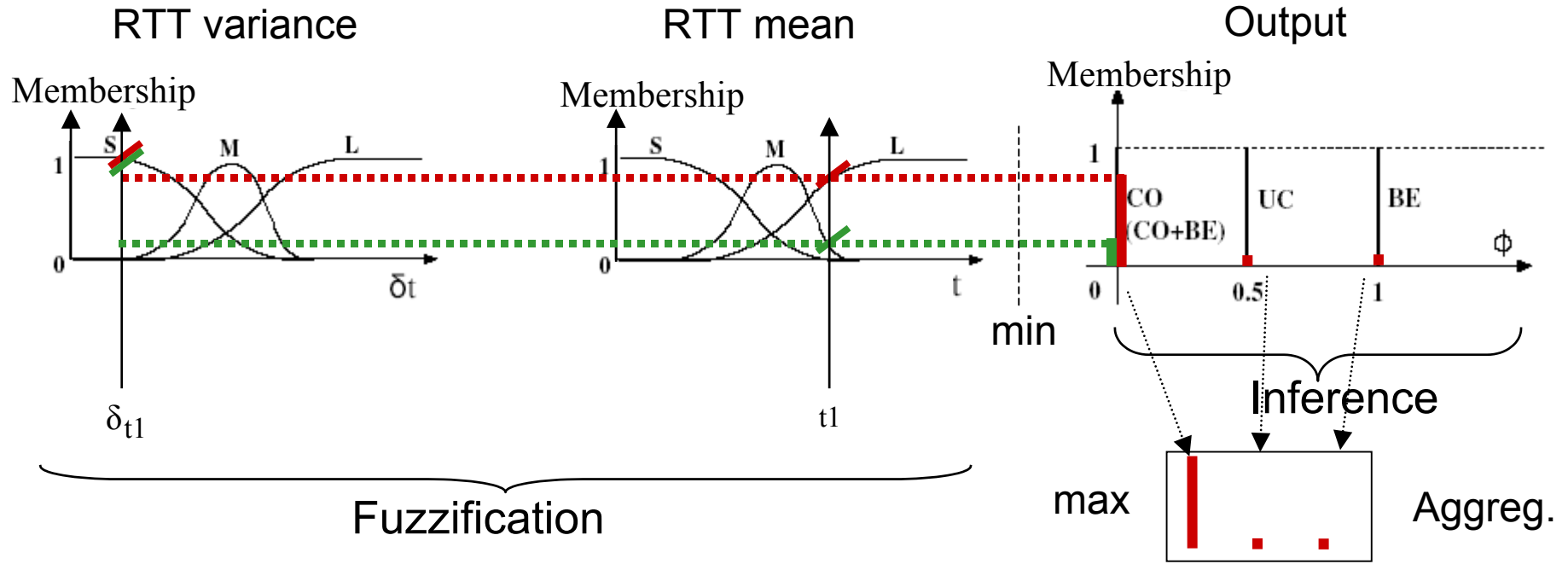
> It uses membership functions and rules to reason about the data

> FL avoids rigid boundaries  
ex. "people height":  
is 1.79m short and 1.80m tall?



> FL is proper for pattern recognition tasks

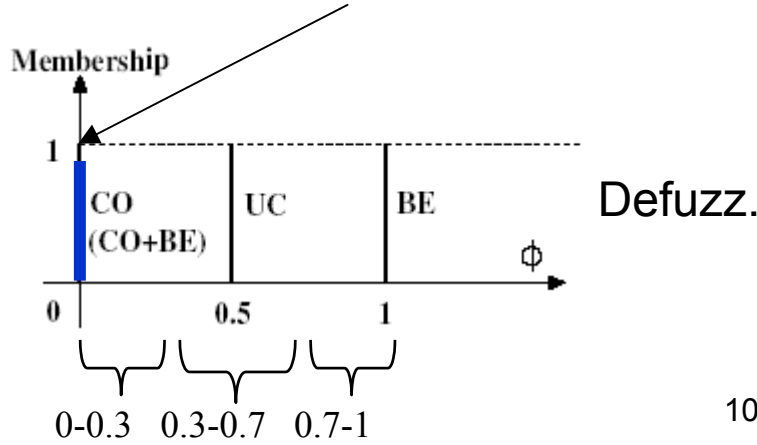
# The Inference Procedure of the Fuzzy Engine



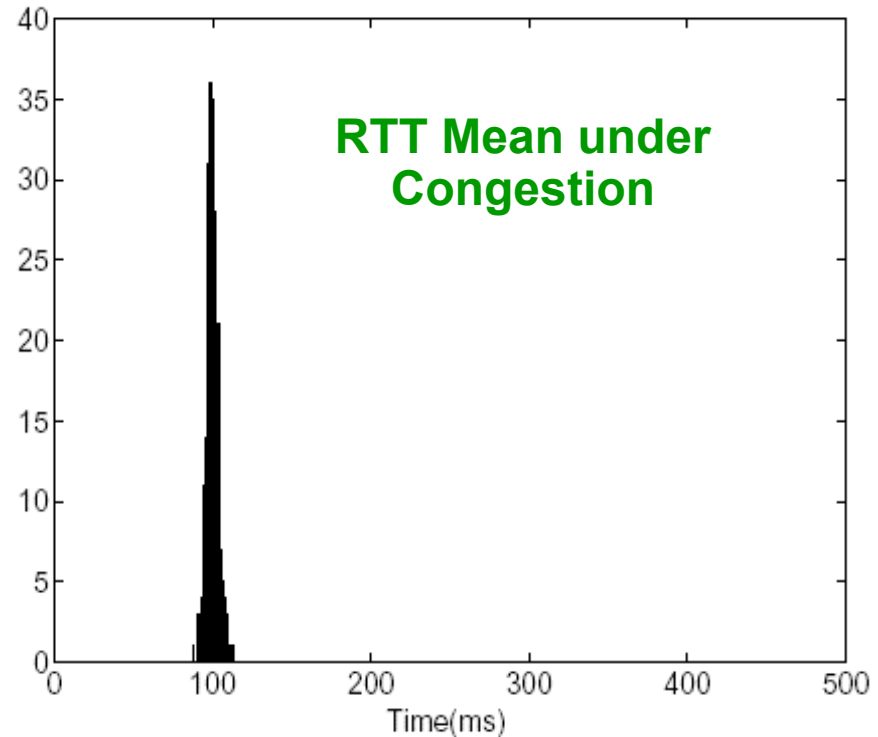
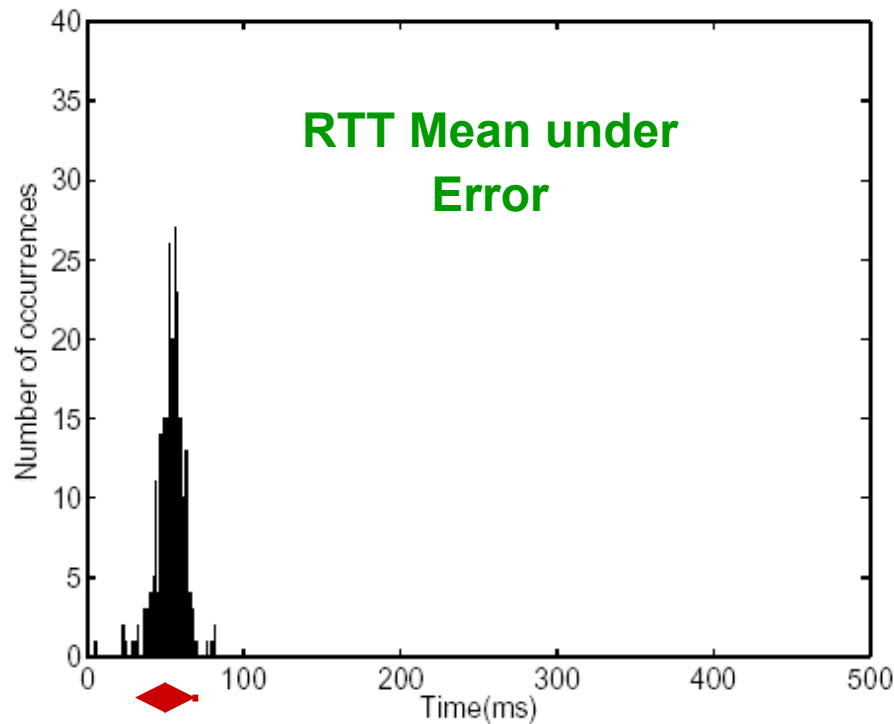
Fuzzy rules output  $\Phi$

| Var( $\delta$ ) \ mean (t) | S  | M  | L  |
|----------------------------|----|----|----|
| S                          | BE | CO | CO |
| M                          | BE | UC | CO |
| L                          | BE | BE | CO |

S = Small  
 M = Medium  
 L = Large  
 CO = Congestior  
 UC = Uncertain  
 BE = Bit Error



# Settings of the Membership Functions for RTT mean by Simulations (3-hop)



Being conservative:  $RTT > 60 \rightarrow$  congestion  
 $RTT < 40 \rightarrow$  medium error

Fuzzy settings: Gaussian curve center is 50 and width is 20

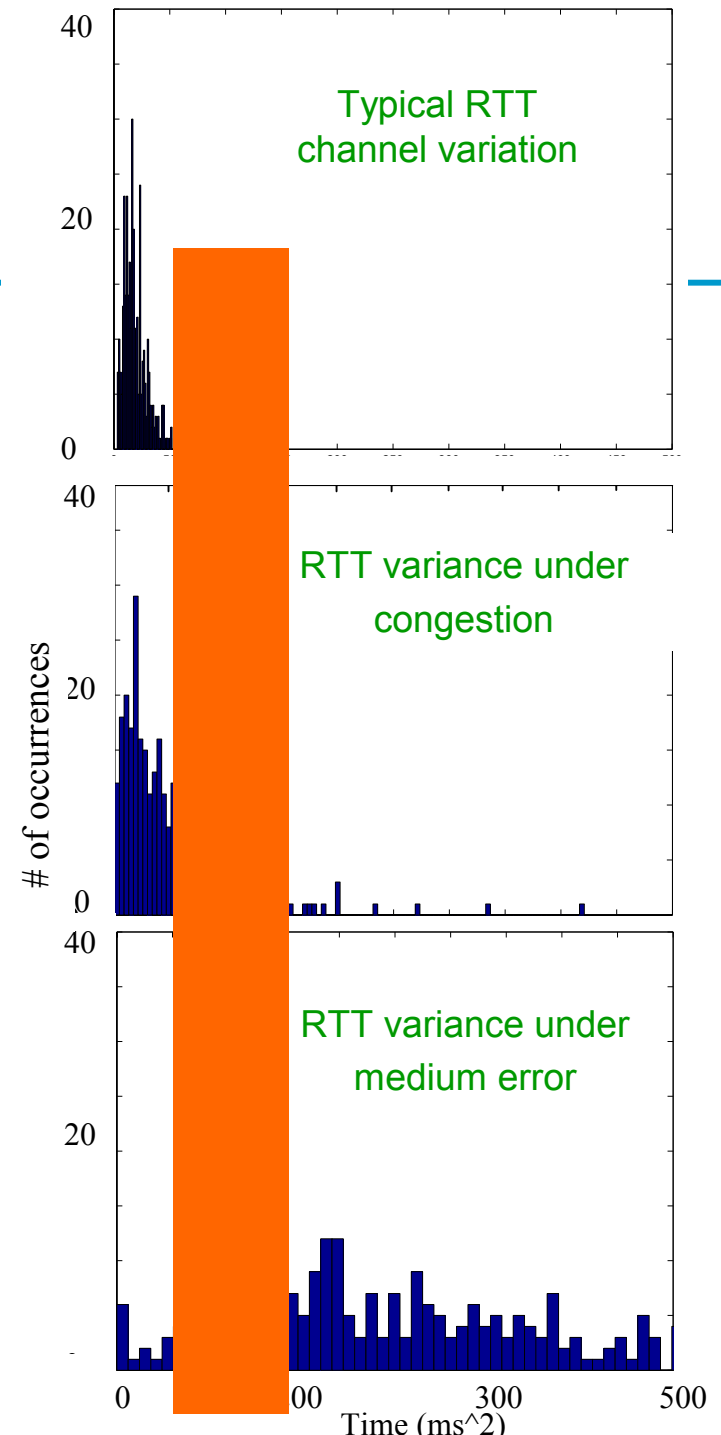
## Settings of the Membership Functions for RTT var. (3-hop)

- > Congestion and medium error have distinct distributions for RTT variance
- > There is an overlapping area between 50-150 ( $\text{ms}^2$ )
- > Below 50 the values are clearly low and above 150 clearly high



Fuzzy settings:

- > Gaussian curve: center is 100 and width is 50.
- > Universe of discourse is 500



# Fuzzy Engine Features

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- > **Fuzzy Input**
  - RTT mean ( $t$ ) and variance ( $\delta$ )

- > **Fuzzy Rules**

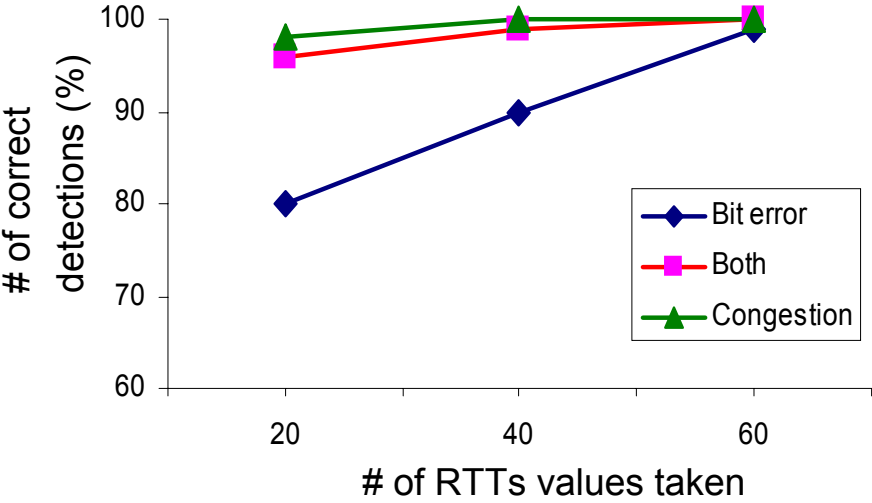
Fuzzy rules output  $\Phi$

| $\delta_t$ | S  | M  | L  |
|------------|----|----|----|
| S          | BE | CO | CO |
| M          | BE | UC | CO |
| L          | BE | BE | CO |

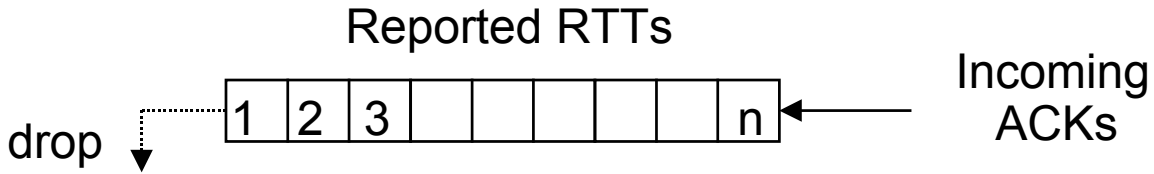
- > **Settings of the membership functions**
  - ✓ mean ( $t$ ): 50 +/- 20 ms
  - ✓ Variance ( $\delta$ ): 100 +/- 50 ms<sup>2</sup>
  - ✓ Universe of discourse is 500 ms<sup>2</sup>
  - ✓ Output ranges:
    - CO: 0.0-0.3
    - UC: 0.3-0.7
    - BE: 0.7-1.0

# Fuzzy engine output correctness

Correct detection (%) vs. # of RTTs per sampling  
*Worst case*



### History of RTTs



## Results

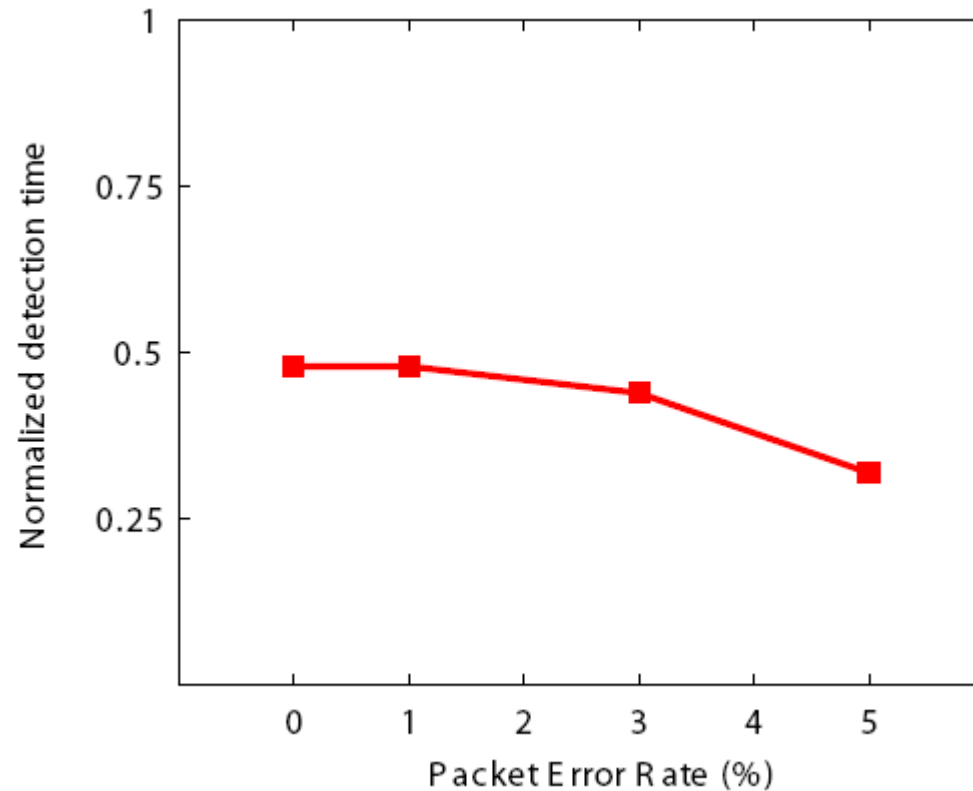
- Congestion was detected in over 98% of the cases
- Bit error requires more values

## scenario

- Chain topology with 3 hops
- 3 distinct runs of 100 seconds
  - Packet error rate of 10%
  - 1 competing flow for cong.
  - Both (PER=5% + 1 c. flow)

# Fuzzy engine reaction time to abrupt congestion

Detection time vs. error rate



- > Scenario: an unconstrained channel is suddenly overloaded
- > The fuzzy engine detects incipient congestion in time

# Summary of the Fuzzy Engine based Approach

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- > If correctly adjusted it may provide accurate results
- > Most effective in steady state conditions
- > It is very computational inexpensive
- > It cannot always obtain a conclusive result
- > Tradeoff between speed and accuracy is always an issue
- > RTT measurements tend to become unfeasible for large number of hops

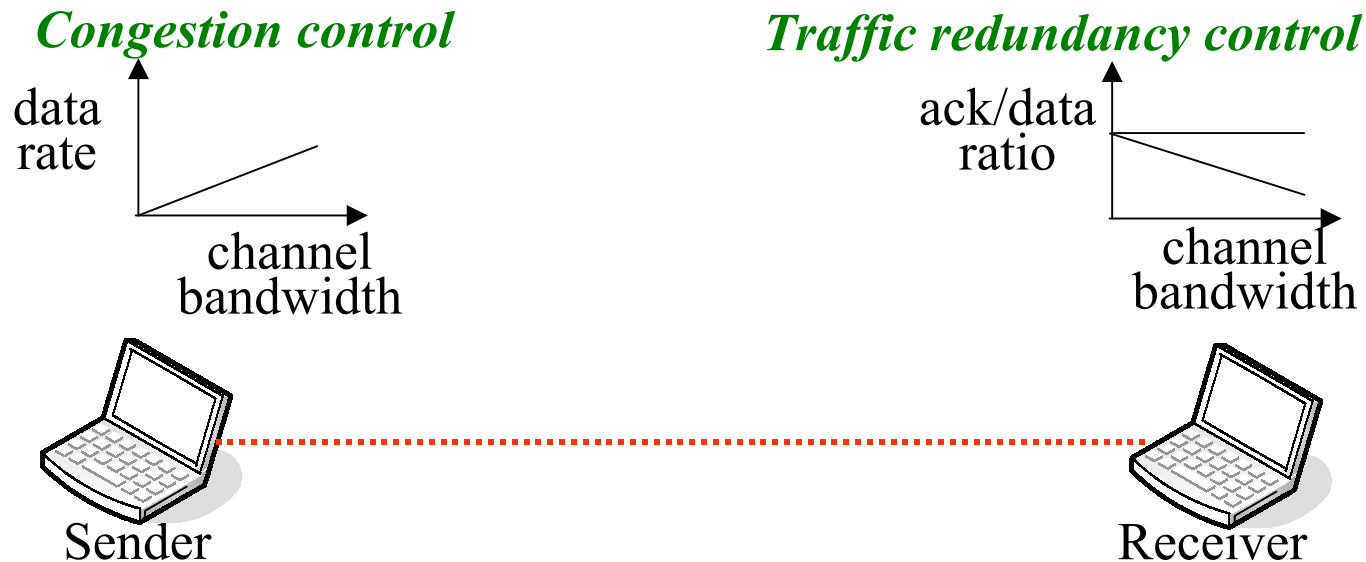
## Future work

- > Different scenarios
- > More elaborated fuzzy engines
- > Adaptive algorithms for setting the fuzzy engine parameters
- > Integration with the “error recovery” mechanism



# Second Main Contribution: A Smart TCP Acknowledgment Approach

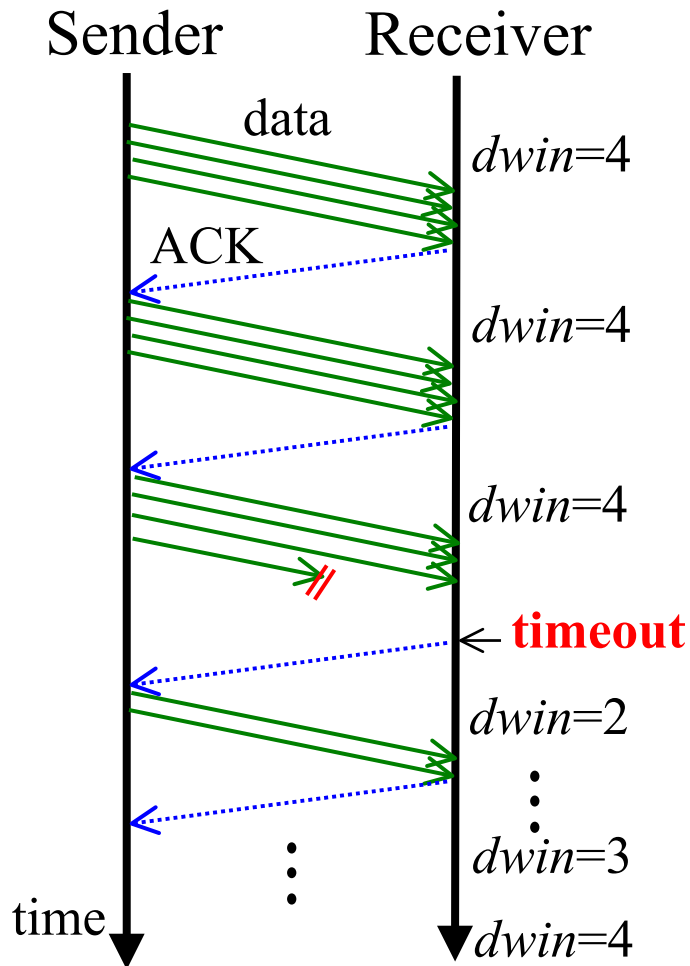
## Smart Acknowledgment



## Dynamic Adaptive Acknowledgements (DAA)

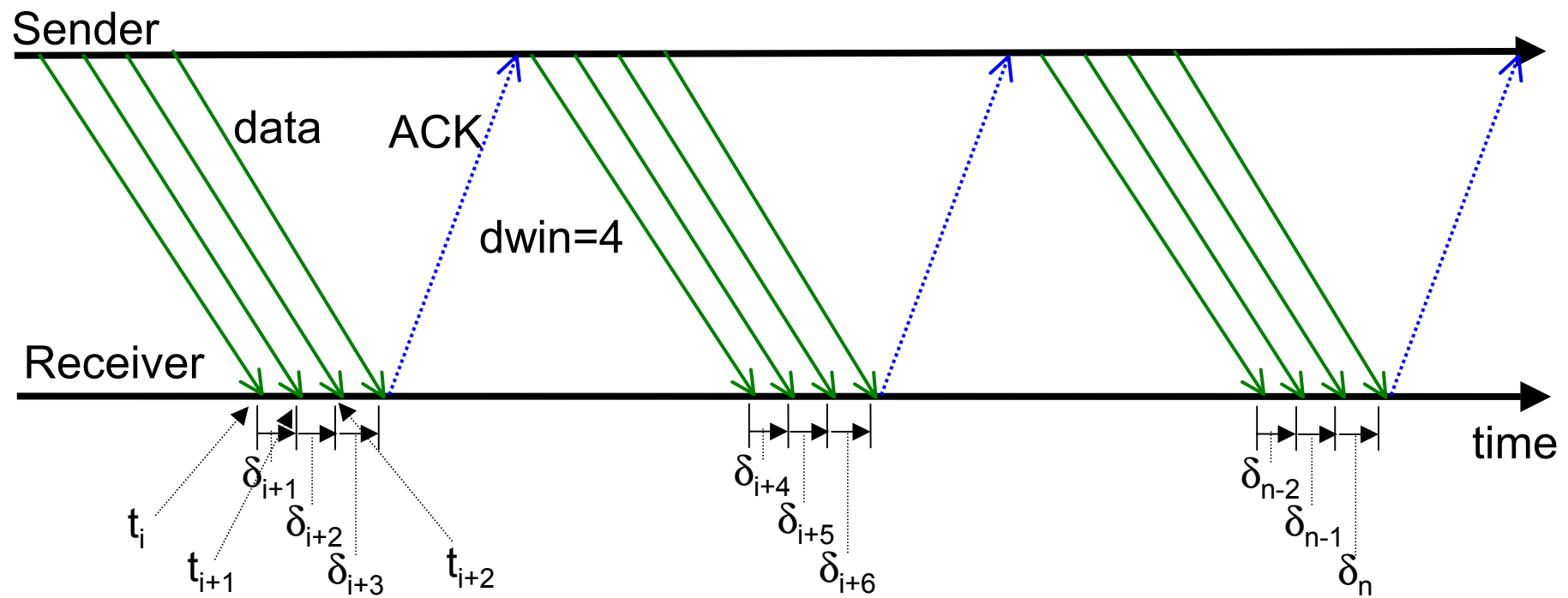
- Less ACKs leave more bandwidth for data packets
- Avoid useless retransmissions ! save energy
- Self-adjustment depending on channel conditions

# DAA Mechanisms at the Receiver



- > Receiver combines up to 4 ACKs ( $dwin = 4$ )
- > Out-of-order packet or timeout triggers immediate ACK !  $dwin := 2$
- >  $dwin$  is increased by 1 (by  $\mu < 1$  if in startup phase)
  - for each correctly received data packet (limit  $dwin = 4$ )
- > Timeout interval  $T$  is adapted based on data packet inter-arrival time  $\delta$

# Packet Inter-arrival computation



At  $t_i, t_{i+1}, t_{i+2}, \dots \rightarrow$  start timer and increment delack counter (ack\_count)

**Low-pass filter:**

$$\bar{\delta}_i = (1-\alpha) \bar{\delta}_{i-1} + \alpha \delta_i$$

$$T_i = (2+k) \bar{\delta}_i$$

- >  $\bar{\delta}_i$  = smoothed pkt inter-arrival
- >  $\bar{\delta}_{i-1}$  = sampled pkt inter-arrival
- >  $\alpha$  = smoothing factor
- >  $k$  = timeout tolerance factor
- >  $T_i$  = timeout interval at arrival of pkt  $i$

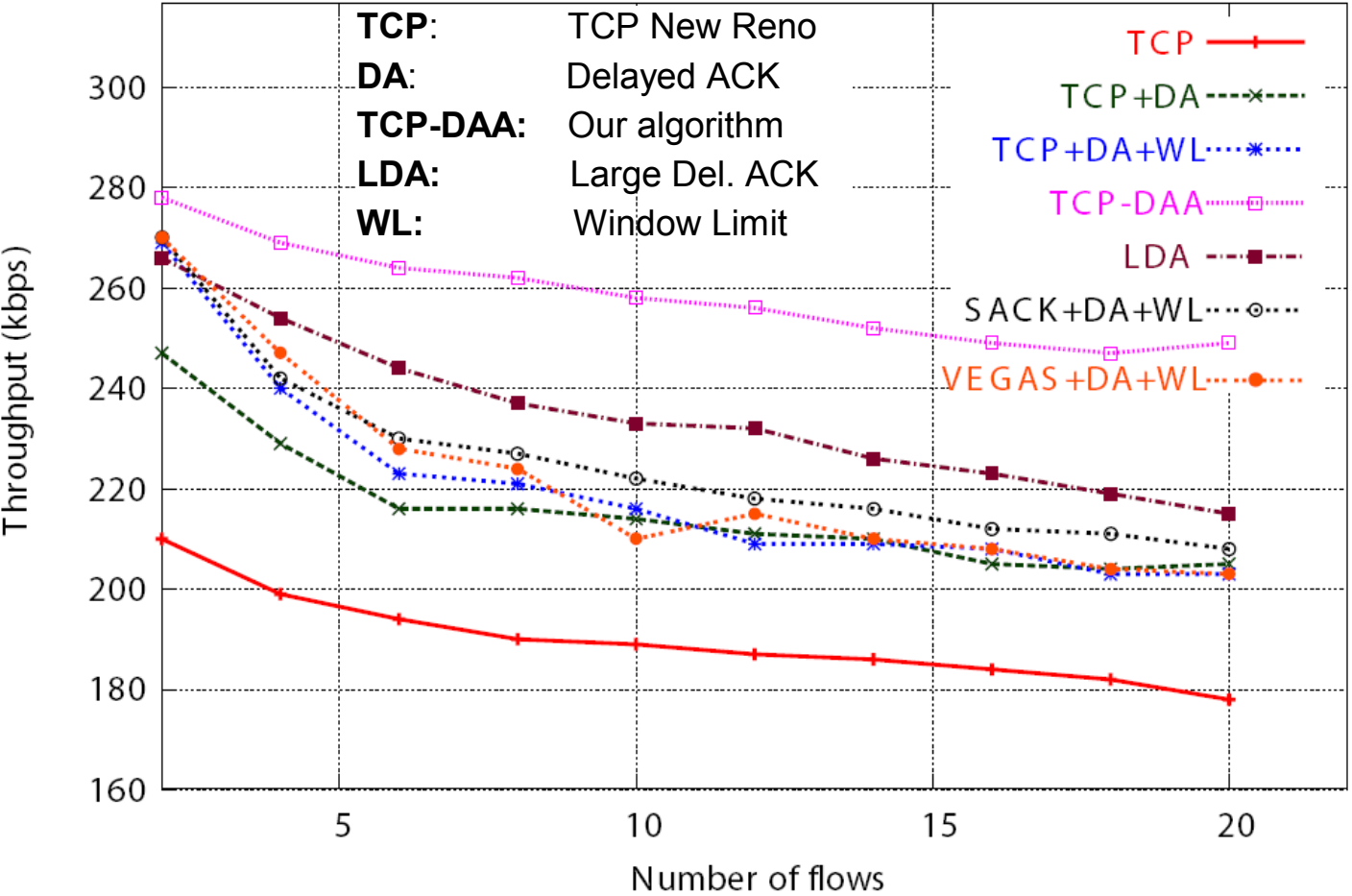
# Adjustments at the DAA Sender

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- > Fivefold retransmission timeout interval ( $\text{newRTO} = 5 * \text{RTO}$ )
- > Fast retransmission for 2 duplicate ACKs
- > Maximum congestion window = 4  
(enough to as many as 10 hops)

# Throughput

TCP throughput vs. network load (hops: 5)

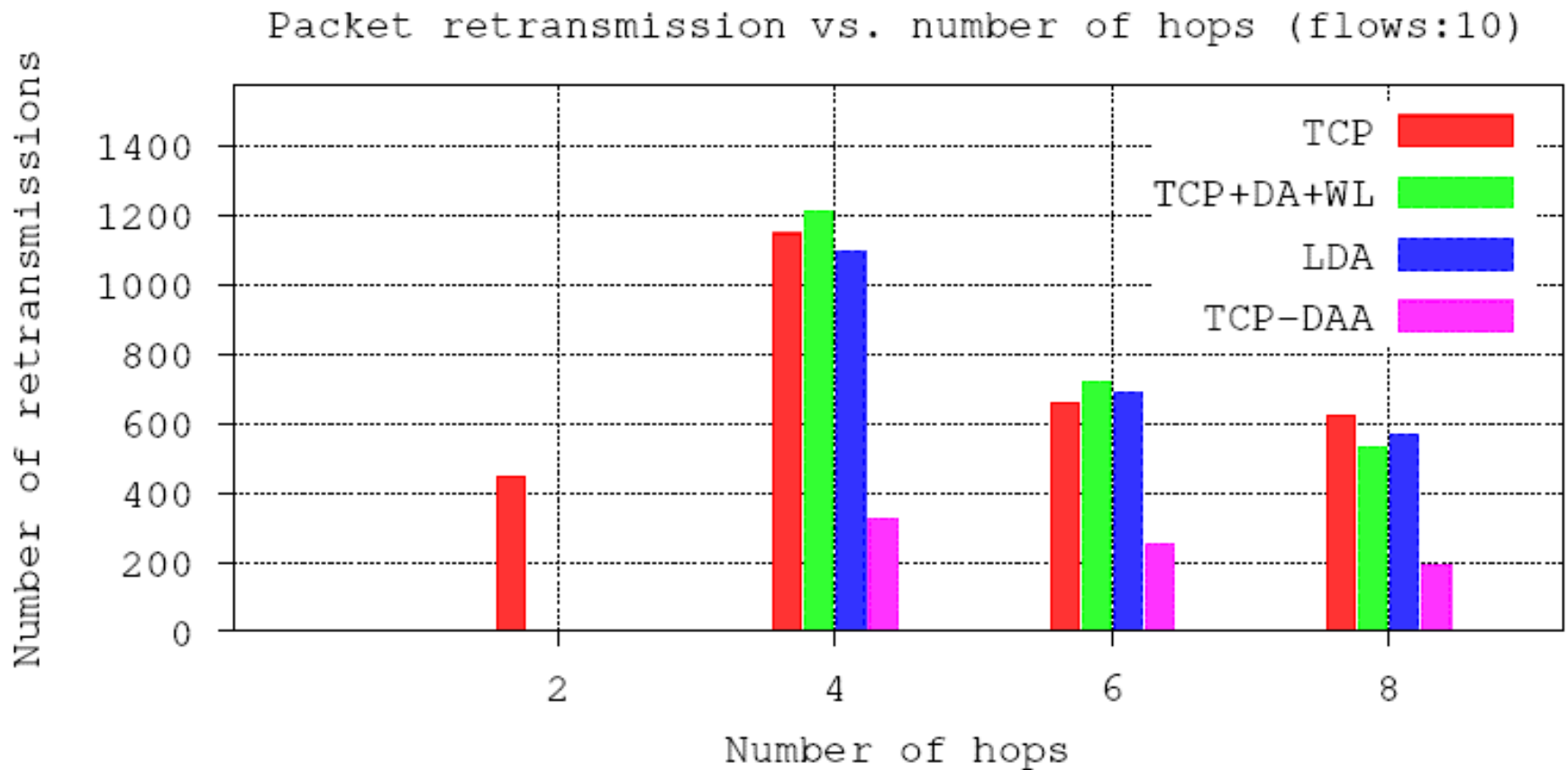


### Simulation scenario

- chain topology
- loss due to collisions by MAC protocol only
- duration: 300 seconds

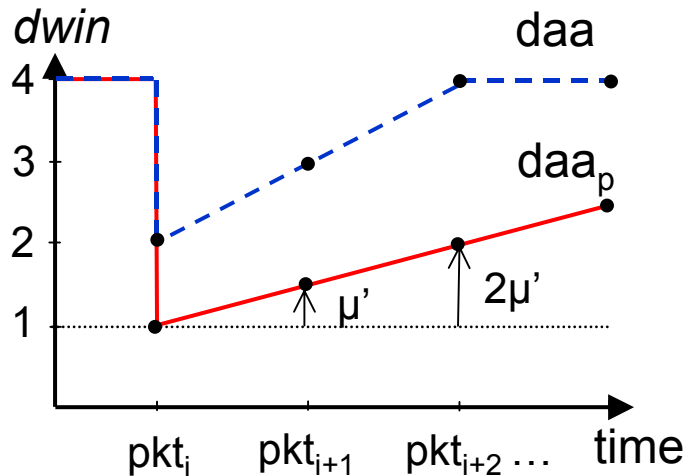
# Retransmissions

- > Lower traffic overhead leads to fewer retransmissions
- > Energy consumption benefits are considerable



# A Modified dwin strategy for Robustness in Highly noisy Environments (DAAp)

- > *dwin* decreases more and increases slower than basic algorithm
  - > Sender regular RTO is doubled
- (W: cumulative value of cwnd)



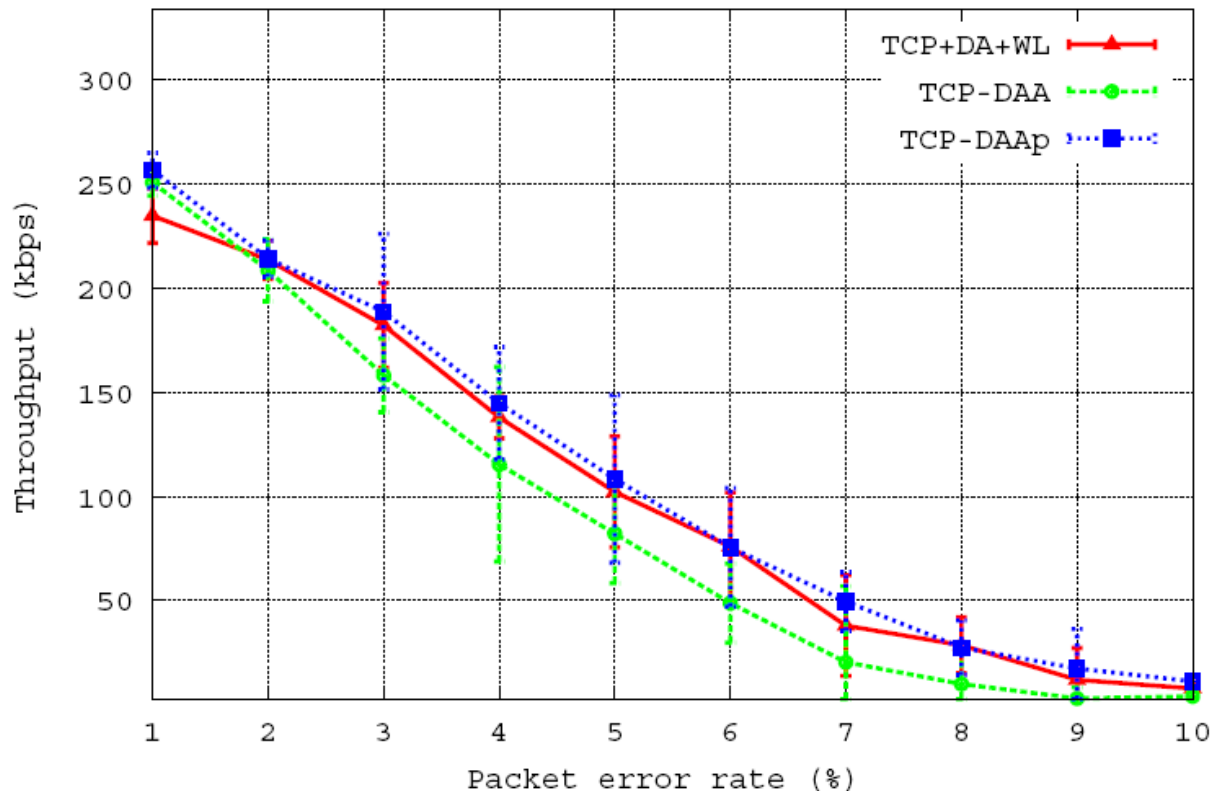
- > 
$$cwnd_i = cwnd_{i-1} + \frac{1}{cwnd_{i-1}}$$
- > 
$$W = cwnd_0 + \sum_{i=1}^n \frac{1}{cwnd_i}$$
- > for  $W=4$  (max cwnd)  $\rightarrow n=7 \rightarrow \mu' = 0.3$

- > 
$$dwin_i = \begin{cases} dwin_{i-1} + \mu, & \text{if in startup} \\ dwin_{i-1} + \mu', & \text{otherwise} \end{cases}$$

# Throughput under external disturbances

## Comparison with adjusted TCP NewReno

TCP throughput vs. error rate (hops: 5, flows: 1)



- > DAAp prevents the sender from missing ACKs
- > DAAp sender reacts faster than DAA
- > Switching mechanism is useful



# Summary of the Smart Acknowledgment Approach

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- > The proposed mechanism improves throughput and energy consumption in a variety of scenarios
- > It outperforms related approaches in many situations
- > It is proper for environments under moderate loss rates

## Future work

- > A tailored sender side algorithm
- > Optimization of the configuration parameters

## Conclusions and outlook

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- > The IEEE 802.11 is feasible for short-range networks
- > Existing protocols can be fine tuned to enhance 802.11 performance
- > Complex modeling are mostly unfeasible here
- > Proactive and reactive approaches should be combined
- > End-to-end approaches facilitate deployment
  
- > *An alternative MAC protocol for long-range ad hoc networks is very much needed*
- > *TCP will have to be adjusted to such a new protocol*
- > *Cross layer design seems to be the way*